

Programs Horn clauses and unification, with backtracking
conclusion :- condition,
condition, condition.

Grammar Term ::= Constant |
CompoundTerm | Variable
Constant ::= Atom | Number
CompoundTerm
::=Term(Term+)
Variables start with capital letters
Atoms start with small letters

Operators Define using op(precedence,
associativity, name).
Built in: , (and), ; (or), is
(arithmetic expression , RHS
must be ground), =:=, =\=
([not] arithmetic equal), <=, >=
(arithmetic smaller / larger), =,
\= ([not] unifiable), ==, \==,
([not] identical)
Note: "is" cannot solve equations

Lists Cons = .(head, tail), Empty= []
Alternate syntax: [a, b, c | Rest]

Cut ! : controls backtracking. When
the cut is "crossed" backtracking
over it will cause the clause to
fail
"fail" always fails, so you can use
!, fail to fail fast
Green cuts prune the search tree
(as in "member", later)
Red cuts omit explicit conditions:
program becomes special cases
and a default rule

Negation not(G) :- call(G), !, fail.
Negation is failure in Prolog, do
not confuse this with logical
negation! For example:
innocent(a).
?- innocent(b). = <No>
Solve by storing negative
information, adding a clause
saying that if it occurs once it
cannot occur anywhere else (e.g.
a battle only occurs in 1815, then
don't both adding "battle didn't
occur in 1816" etc) or make the
closed world assumption
(everything that exists is stated
in the program) and implement
negation by failure.

Extras _ represents unused variable
Declare rule with "dynamic
ruleName\arity" e.g. "dynamic

awayGoal\2"
Use assert/retract to introduce
clauses into Prolog at runtime

Basic Examples length([], 0).
length([H|T], N) :- length(T, NT),
N is NT + 1.
length2(L, N) :- acc(L, 0, N).
acc([], A, A).
acc([H|T], A, N) :- A1 is A + 1,
acc(T, A1, N). (tail recursive)
max([], A, A).
max([H|T], A, M) :- H > A,
max(T, H, M).
max([H|T], A, M) :- H =< A,
max(T, A, M). (initialise the
accumulator with the list head)
reverse([], []).
reverse([H|T], O) :- reverse(T,
Y), append(Y, [H], O).
reverse2([], O, O).
reverse2([H|T], A, O) :-
reverse2(T, [H|A], O). (tail rec.)

Partial Maps partial([], []).
partial([X|T], [X|L]) :- include(X),
partial(T, L).
partial([X|T], L) :- partial(T, L).
partial2([], [], []).
partial2([H|T], [H|A], B) :-
inA(H), partial2(T, A, B).
partial2([H|T], A, [H|B]) :-
partial2(T, A, B).

Graphs Binary Search Tree Represent edges as a(u,v) or u-v
mt. % empty node
n(L,Value,R). % node functor

insert(Item, n(L, Item, R), n(L,
Item, R)).
insert(Item, mt, n(L, Item, R)).
insert(Item, n(L, T, R), n(NL, T,
R)) :- Item < T, insert(Item, L,
NL).
insert(Item, n(L, T, R), n(L, T,
NR)) :- Item > T, insert(Item, R,
NR).

Using Cut member(X, [X|_]) :- !.
member(X, [_|L]) :- member(X,
L).
Warning:
max(X, Y, X) :- X >= Y, !.
max(X, Y, Y).
max(10, 0, 0) = Yes
max(10, 0, X) = 10, <No>

Difference Lists

Use:

$\text{max}(X, Y, X) :- X \geq Y, !.$

$\text{max}(X, Y, Y) :- X < Y, !.$

X-X is an empty list

Change list X to a normal list by unifying it with Y-[]

$\text{append}(A-B, B-C, A-C).$ (constant time append: execute as

$\text{append}([a, b|Y]-Y, [c, d|Z]-Z, X-XS).$)

$\text{reverse}([]-[], []-[]).$

$\text{reverse}([H1|T]-\text{End1}, \text{Result}) :-$

$\text{reverse}(T-\text{End1}, \text{SubResult}),$

$\text{append}(\text{SubResult}, [H1|X]-X, \text{Result}).$

Better, not using append:

$\text{reverse}([]-[], X-X).$

$\text{reverse}([H|T]-\text{End1}, \text{Start2}-$

$\text{End2}) :- \text{reverse}(T-\text{End1}, \text{Start2}-[\text{Head}|\text{End2}]).$