| Definitions IEEE Floating Point $\pm 0: e, m = 0$ $\pm \infty: e = max, m = 0$ NaN: $e = max, m \neq 0$ Denorm: $e = 0, m \neq 0$ | | Directories FS Issues | Directories stored as files open/create (SFID -> UFID) | | | |
|--|--|--|--|--|--|--|
| MultiprogrammingSchedulingFirst come, first servedShortest job first (optimal)Shortest remaining time first (new | | | | Concurrency control: locks. Mandatory/advisory, shared/exclusive | | |
| l e Average Priority | process w Round rot $	au_{n+1} = at_n +$ Static: e.g | / low burst: preempt) bin (with quantum) $-(1-a)\tau_n$ J. use priority as next | Protectic Goals | Prevent information disclosure/modification Denial of service | | |
| i | Dynamic: | burst time aging (prevent starving) ted: penalise time use | Mechanisr | Isolation (debug/error control) ns User/supervisor modes Memory management control File control (ACLs etc) | | |
| Memory Memory | Dynami | ixed size partitions c: partition at runtime physical frame to logical | Principles | Physical restrictions Passwords/encryption Stupidity/legislation Least privilege Default deny | | |
| Compaction | Run tim | e relocation hen you move off swap | | Current authority (caching) Psychologically acceptable | | |
| Replacement | FIFO: B LRU: tin | elady's anomaly etc nestamp/page stack ference/dirty bit or | Authentica | High circumvention cost | | |
| Segments | second Referen Page bu App-spe Locality "Workin Local/gl Externa Softwar | chance FIFO (clock) ce counting iffering (pool of victims) ecific hooks of reference g set" <=> thrashing obal page tables I fragmentation e segments (page array, | Access Ma Capabilitie | rix Keyed on object -> ACL Keyed on subject -> capability | | |
| | OS keeps priv. consistent) Paged segments (per-segment | | | Unix File System | | |
| I/O | - | bles, not portable) | Inodes Directory | Type/mode/user groupid/size/nlinks Direct x12/single doub trip indirect Files with inodes holding list of SFID | | |
| Access Modes | s Polled | vs interrupt driven | | Can have at most 1 hard link | | |
| Buffering | Mainta | ng/nonblocking/asynch ins copy semantics 'double/circular | Disk | Boot super inode table data blocks Superblock: nfree, free link lists etc Can "mount" into name service | | |
| Other issues | Sized a Cachin (queue | Sized according to device type Files Descriptor tabl Caching, scheduling System wide -> (queue/fairness), device UGO bits + set | | Descriptor table: process specific -> system wide -> device inode table UGO bits + setg uid. Directories use X = cwd, SG = group "sticky" | | |
| File Issues | Directo Storag | ory service (name -> id) e service (id -> data) S must be implemented | | Consistency issues (on crash) | | |

Unix Processes

| Principles | Heavyweight (own page table, are the unit of scheduling) |
|------------|--|
| | Shared kernel space -> no c-switch Zombie state (for parents benefit) |
| Boot | Kernel -> init -> tty -> login -> sh |
| IPC | Pipes (later named pipes): consist of finite circular queue |
| 1/0 | Signals which process can catch |
| I/O | Buffer cache w/ sync every 30 seconds |
| Cabadulina | Aggressive metadata writeback |
| Scheduling | Lower priorities superuser only Penalises CPU usage over $\approx 5s$ |

Windows Architecture

| Structure | Super: HAL, kernel, executive |
|------------|--|
| | User: environment/protection subs. |
| | |
| | HAL: interrupt/DMA/SMP etc |
| | Kernel: no pre-emption, schedules, |
| Dragona | handles interrupts, processor sync. |
| Processes | Processes own resources |
| | Threads are dispatch units, |
| | lightweight and share resources |
| | Parent/child not mandatory |
| Scheduling | Boost on return from IO/fg thread |
| | Priority decays over time to base |
| | Also get static priority ("real time") |
| Objects | Object manager checks |
| | ACLs/creates objects handles |
| | Implies uniform security model |
| | enforced by Security Ref. Manager |
| | Name, directory, security |
| | descriptor, type info, ref count |
| | Live in a namespace with recursive |
| | name parsing responsibility |
| VMM | Can share memory in section |
| | object = segment (based non) |
| IPC | Channels (copy, zero-copy, quick) |
| I/O | Asynchronous: IRP holds |
| | parameters, results etc |
| | Stackable drivers handle IRPs |
| | "Virtual block" cache w/ prefetch |
| | Unified cache works on VAS "lines" |
| | (VMM does cache I/O) |
| | User control (temp/write through) |
| Subsystems | Layered over NT native API |
| | DOS, OS/2, POSIX, WoW |

Windows File Systems

| FAT16 | Linked list of clusters: max 2Gb |
|-------|----------------------------------|
| | Variable cluster size |
| FAT32 | Wider FAT16: 8Gb @ 4k cluster |

NTFS

Root directory anywhere on disk Can use backup FAT (fault tolerant) VFAT: long names on top of this File records held in MFT (itself a file) indexed on file ref (64 bits) Based on a volume, not partition Files are attribute/value pairs Special: LogFile, Bitmap, BadClus Transactions for consistency Volumes can be RAID sets, supporting bad cluster remapping Security descriptor in MFT Compression and sparsity Symmetric encryption w/ RSA key on that key, admin can get key..